

Relational Operations: Overriding

$$r = \{(a, 1), (b, 2), (c, 3), (a, 4), (b, 5), (c, 6), (d, 1), (e, 2), (f, 3)\}$$

Example: Calculate r overridden with $\{(a, 3), (c, 4)\}$

Hint: Decompose results to those in t 's domain and those not in t 's domain.

$$\begin{aligned} \underbrace{(r)}_{\text{relation}} \triangleleft \underbrace{(t)}_{\text{relation}} &= \{(d, r') \mid (d, r') \in t \vee \underline{((d, r') \in r \wedge d \notin \text{dom}(t))}\} \\ &= \{(d, r') \mid (d, r') \in t\} \cup \{(d, r') \mid (d, r') \in r \wedge d \notin \text{dom}(t)\} \end{aligned}$$

Exercises: Algebraic Properties of Relational Operations

$$r = \{(a, 1), (b, 2), (c, 3), (a, 4), (b, 5), (c, 6), (d, 1), (e, 2), (f, 3)\}$$

Define the **image** of set s on r in terms of other relational operations.

Hint: What range of value should be included?

Define r **overridden with** set t in terms of other relational operations.

Hint: To be in t 's domain or not to be in t 's domain?

Functional Property

$\text{isFunctional}(r) \Leftrightarrow$

$\forall s, t1, t2 \bullet$

$(s \in S \wedge t1 \in T \wedge t2 \in T)$

\Rightarrow

$((s, t1) \in r \wedge (s, t2) \in r \Rightarrow t1 = t2)$

Q: Smallest relation satisfying the functional property.

Q: How to **prove** or **disprove** that a relation r is a function.

Q: Rewrite the functional property using **contrapositive**.

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Partial Functions vs. Total Functions

$$r \in S \rightharpoonup T \Leftrightarrow (\text{isFunctional}(r) \wedge \text{dom}(r) \subseteq S)$$

$$r \in S \rightarrow T \Leftrightarrow (\text{isFunctional}(r) \wedge \text{dom}(r) = S)$$

Exercise: Visualize $S \rightharpoonup T$ vs. $S \rightarrow T$

e.g., $\{ \{(2, a), (1, b)\}, \{(2, a), (3, a), (1, b)\} \} \subseteq \{1, 2, 3\} \rightharpoonup \{a, b\}$

e.g., $\{(2, a), (3, a), (1, b)\} \in \{1, 2, 3\} \rightarrow \{a, b\}$

e.g., $\{(2, a), (1, b)\} \notin \{1, 2, 3\} \rightarrow \{a, b\}$

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Relational **Image** vs. Functional **Application**

A function is a relation.

$$f \in \{1, 2, 3\} \rightarrow \{a, b\}$$

$$f = \{ (3, a), (1, b) \}$$

Exercises:

$$f[\{3\}] =$$

$$f[\{1\}] =$$

$$f[\{2\}] =$$

Modelling Decision: Relations vs. Functions

An organization has a system for keeping track of its employees as to where they are on the premises (e.g., ``Zone A, Floor 23``). To achieve this, each employee is issued with an active badge which, when scanned, synchronizes their current positions to a central database.

Assume the following two sets:

- *Employee* denotes the **set** of all employees working for the organization.
- *Location* denotes the **set** of all valid locations in the organization.

Is $\text{where_is} \in \text{Employee} \leftrightarrow \text{Location}$ appropriate?

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Is $\text{where_is} \in \text{Employee} \nrightarrow \text{Location}$ appropriate?

Injective Functions

$isInjective(f)$

\iff

$$\forall s_1, s_2, t \bullet (s_1 \in S \wedge s_2 \in S \wedge t \in T) \Rightarrow ((s_1, t) \in f \wedge (s_2, t) \in f \Rightarrow s_1 = s_2)$$

If f is a **partial injection**, we write: $f \in S \rightsquigarrow T$

- e.g., $\{\emptyset, \{(1, \mathbf{a})\}, \{(2, \mathbf{a}), (3, \mathbf{b})\}\} \subseteq \{1, 2, 3\} \rightsquigarrow \{a, b\}$
- e.g., $\{(1, \mathbf{b}), (2, a), (3, \mathbf{b})\} \notin \{1, 2, 3\} \rightsquigarrow \{a, b\}$
- e.g., $\{(1, \mathbf{b}), (3, \mathbf{b})\} \notin \{1, 2, 3\} \rightsquigarrow \{a, b\}$

If f is a **total injection**, we write: $f \in S \succrightarrow T$

- e.g., $\{1, 2, 3\} \succrightarrow \{a, b\} = \emptyset$
- e.g., $\{(2, d), (1, a), (3, c)\} \in \{1, 2, 3\} \succrightarrow \{a, b, c, d\}$
- e.g., $\{(2, d), (1, c)\} \notin \{1, 2, 3\} \succrightarrow \{a, b, c, d\}$
- e.g., $\{(2, \mathbf{d}), (1, c), (3, \mathbf{d})\} \notin \{1, 2, 3\} \succrightarrow \{a, b, c, d\}$

Surjective Functions

$$isSurjective(f) \iff \underline{ran}(f) = \underline{T}$$

If f is a **partial surjection**, we write: $f \in S \twoheadrightarrow T$

- e.g., $\{ \{(1, \mathbf{b}), (2, \mathbf{a})\}, \{(1, \mathbf{b}), (2, \mathbf{a}), (3, \mathbf{b})\} \} \subseteq \{1, 2, 3\} \twoheadrightarrow \{a, b\}$
- e.g., $\{(2, \mathbf{a}), (1, \mathbf{a}), (3, \mathbf{a})\} \notin \{1, 2, 3\} \twoheadrightarrow \{a, b\}$
- e.g., $\{(2, \mathbf{b}), (1, \mathbf{b})\} \notin \{1, 2, 3\} \twoheadrightarrow \{a, b\}$

If f is a **total surjection**, we write: $f \in S \rightarrow T$

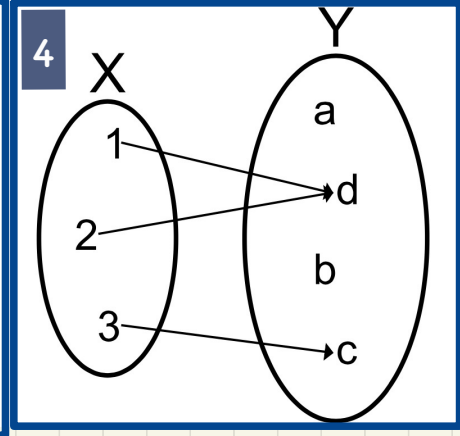
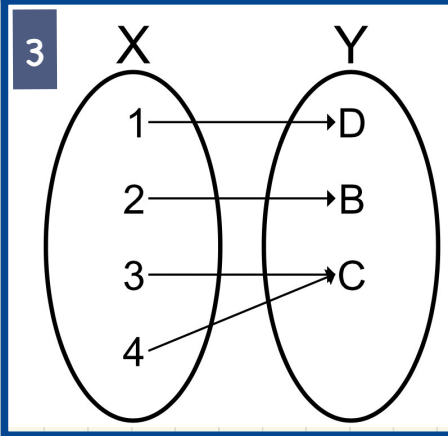
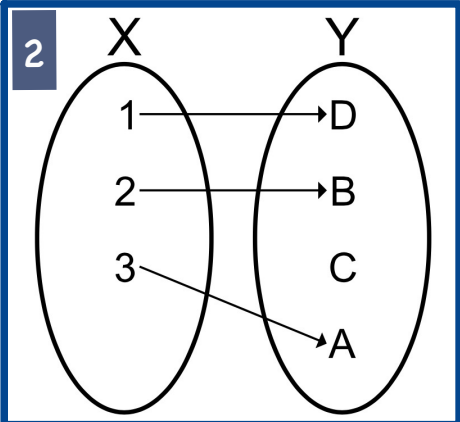
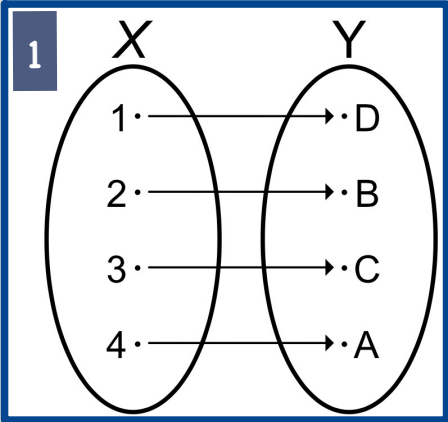
- e.g., $\{ \{(2, a), (1, b), (3, a)\}, \{(2, b), (1, a), (3, b)\} \} \subseteq \{1, 2, 3\} \rightarrow \{a, b\}$
- e.g., $\{(2, \mathbf{a}), (3, \mathbf{b})\} \notin \{1, 2, 3\} \rightarrow \{a, b\}$
- e.g., $\{(2, \mathbf{a}), (3, \mathbf{a}), (1, \mathbf{a})\} \notin \{1, 2, 3\} \rightarrow \{a, b\}$

Bijjective Functions

f is *bijjective/a bijection/one-to-one correspondence* if f is *total*, *injective*, and *surjective*.

- e.g., $\{1, 2, 3\} \twoheadrightarrow \{a, b\} = \emptyset$
- e.g., $\{ \{(1, a), (2, b), (3, c)\}, \{(2, a), (3, b), (1, c)\} \} \subseteq \{1, 2, 3\} \twoheadrightarrow \{a, b, c\}$
- e.g., $\{(2, b), (3, c), (4, a)\} \notin \{1, 2, 3, 4\} \twoheadrightarrow \{a, b, c\}$
- e.g., $\{(1, a), (2, b), (3, c), (4, a)\} \notin \{1, 2, 3, 4\} \twoheadrightarrow \{a, b, c\}$
- e.g., $\{(1, a), (2, c)\} \notin \{1, 2\} \twoheadrightarrow \{a, b, c\}$

Exercise



	1	2	3	4
partial				
total				
injection				
surjection				
bijection				

Formalizing **Arrays** as **Functions**

```
String[] names = {"alan", "mark", "tom"};
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